

# Microeconomics II - Winter 2005

## Chapter 1

### Games in Strategic Form - Nash equilibrium











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# Summary

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- Preliminaries  
- Dominated Strategies  
- Nash equilibrium: definition  
- Nash equilibrium: examples  
- Nash equilibrium: existence  

## Definition of a game:

- A set of players:  $P = \{1, 2, \dots, I\}$ . A generic player  $i \in P$ , (all others  $-i$ ).
- A set of strategies:  $S_i$ . A generic strategy  $s_i \in S_i$ .  $S = \prod_{i=1}^I S_i$
- Payoff functions for each player:  $u_i : S \rightarrow \mathfrak{R}$ . We write  $u_i(s) = u_i(s_1, \dots, s_I) = u_i(s_i, s_{-i})$ .

Examples:

**A**  $P = \{1, \dots, 18\}$ ,  $S_i = \mathbb{R}^+$ ,  $u_i(s) = 2 \sum_{j=1}^{18} \frac{s_j}{18} - s_i$

**B**  $P = \{1, \dots, 18\}$ ,  $S_i = \mathbb{R}^+$ ,  $u_i(s) = 2 \min_{j \in P} s_j - s_i$

**C**

$sp, bp$	$P$	$N$
$P$	1, 3	-1, 6
$N$	4, 1	0, 0

Size of resource: 6, cost of P:1.

## Mixed strategies:

A mixed strategy for agent  $i$  is a probability distribution over  $S_i$ . That is:

$$\Sigma_i = \left\{ \sigma_i \in \mathbb{R}^{\#S_i} \mid \sigma_i(s_j) \geq 0, \sum_{i \in S_i} \sigma_i(s_i) = 1 \right\}$$

## Payoffs with mixed strategies:

$$\begin{aligned} u_i(\sigma) &= \sum_{s_1 \in S_1} \dots \sum_{s_I \in S_I} \left( \prod_{j=1}^I \sigma_j(s_j) \right) u_i(s) \\ &= \sum_{s_i \in S_i} \sigma_i(s_i) \left( \sum_{s_{-i} \in S_{-i}} \left( \prod_{j=1}^I \sigma_j(s_j) \right) u_i(s_i, s_{-i}) \right) \\ &= \sum_{s_i \in S_i} \sigma_i(s_i) u_i(s_i, \sigma_{-i}) \end{aligned}$$

So payoffs are linear in own strategy and continuous in all strategies.

Example:  $\sigma_{sp} = \left(\frac{1}{3}, \frac{2}{3}\right), \sigma_{bp} = \left(\frac{3}{4}, \frac{1}{4}\right)$

$$\begin{aligned}u_{sp}(\sigma_{sp}, \sigma_{bp}) &= \frac{1}{3} \cdot \frac{3}{4} \cdot 1 + \frac{1}{3} \cdot \frac{1}{4} \cdot (-1) + \frac{2}{3} \cdot \frac{3}{4} \cdot 4 + \frac{2}{3} \cdot \frac{1}{4} \cdot 0 \\&= \frac{1}{3} \left( \frac{3}{4} \cdot 1 + \frac{1}{4} \cdot (-1) \right) + \frac{2}{3} \left( \frac{3}{4} \cdot 4 + \frac{1}{4} \cdot 0 \right) \\&= \frac{1}{3} \cdot \frac{2}{4} + \frac{2}{3} \cdot 3\end{aligned}$$

# Dominated Strategies (1/2)



**A**  $s_i \in S_i$  is **strictly dominated** if  $\exists \sigma_i \in \Sigma_i$  such that

$$u_i(\sigma_i, s_{-i}) > u_i(s_i, s_{-i}) \quad \forall s_{-i} \in S_{-i}$$

This definition is equivalent if we substitute  $s_{-i}$  by  $\sigma_{-i}$ , why?

**B**  $s_i \in S_i$  is **weakly dominated** if  $\exists \sigma_i \in \Sigma_i$  such that

$$u_i(\sigma_i, s_{-i}) \geq u_i(s_i, s_{-i}) \quad \forall s_{-i} \in S_{-i}$$

$$u_i(\sigma_i, s_{-i}) > u_i(s_i, s_{-i}) \quad \text{for some } s_{-i} \in S_{-i}$$

Example: All strategies except 0 are strictly dominated in game A, and  $P$  is strictly dominated for  $sp$ .

## Iterative domination:

Let  $S_i^0 = S_i$  and  $\Sigma_i^0 = \Sigma_i$ . Then, for  $q \geq 1$

$$S_i^q = \left\{ s_i \in S_i^{q-1} \mid \nexists \sigma_i \in \Sigma_i^{q-1} \text{ such that } u_i(\sigma_i, s_{-i}) > u_i(s_i, s_{-i}), \forall s_{-i} \in S_{-i}^{q-1}, \right\}$$

$$\Sigma_i^q = \left\{ \sigma_i \in \Sigma_i^{q-1} \mid \sigma_i(s_i) > 0 \Rightarrow s_i \in S_i^q \right\}$$



# Nash equilibrium: definition (1/2)



A strategy profile  $s^*$  is a *Nash equilibrium* if:

$$u_i(s_i^*, s_{-i}^*) \geq u_i(s_i, s_{-i}^*) \forall s_i \in S_i$$

A strategy profile  $\sigma^*$  is a *Nash equilibrium in mixed strategies* if:

$$u_i(\sigma_i^*, \sigma_{-i}^*) \geq u_i(\sigma_i, \sigma_{-i}^*) \forall \sigma_i \in \Sigma_i$$

Notice here that the definition above is equivalent to:

$$u_i(\sigma_i^*, \sigma_{-i}^*) \geq \sum_{s_i \in S_i} \sigma_i(s_i) u_i(s_i, \sigma_{-i}^*) \forall \sigma_i \in \Sigma_i$$

thus to:

$$u_i(\sigma_i^*, \sigma_{-i}^*) \geq u_i(s_i, \sigma_{-i}^*) \forall s_i \in S_i$$

## Nash equilibrium: definition (2/2)



**Proposition 1** *All strategies in the support of  $\sigma_i^*$  give the same payoff.*

**Proof.** *Suppose not. Then there are  $\sigma_i^*(s'_i)$  and  $\sigma_i^*(s''_i)$  with*

$$u_i(s'_i, \sigma_{-i}^*) > u_i(s''_i, \sigma_{-i}^*)$$

*Then let  $\sigma_i^{**}$  such that  $\sigma_i^{**}(s'_i) = \sigma_i^*(s'_i) + \sigma_i^*(s''_i)$ ,  $\sigma_i^{**}(s''_i) = 0$  and  $\sigma_i^{**}(s_i) = \sigma_i^*(s_i)$  for  $s_i \neq s'_i, s_i \neq s''_i$ . Then we must have  $u_i(\sigma_i^{**}, \sigma_{-i}^*) > u_i(\sigma_i^*, \sigma_{-i}^*)$ , thus a contradiction. ■*

# Nash equilibrium: examples (1/5)



Example 1: Game B. For all  $r \in \mathfrak{R}$ ,  $s = (r, r, \dots, r)$  is a Nash equilibrium.

Example 2:

1,2	L	M	R
T	7,2	2,7	3,6
B	2,7	7,2	4,5

1.(a) No pure strategy equilibrium.

(b) No mixed strategy equilibrium where player 1 uses only pure strategies.

(c) No mixed strategy equilibrium where player 2 uses only pure strategies.



(d) No mixed strategy equilibrium where 1 uses T and B and 2 uses L, M and R.

For this we would need:

$$7\sigma_2(L) + 2\sigma_2(M) + 3(1 - \sigma_2(L) - \sigma_2(M)) = \\ 2\sigma_2(L) + 7\sigma_2(M) + 4(1 - \sigma_2(L) - \sigma_2(M))$$

and

$$2\sigma_1(T) + 7(1 - \sigma_1(T)) = 7\sigma_1(T) + 2(1 - \sigma_1(T)) = 6\sigma_1(T) + 5(1 - \sigma_1(T))$$

But the first of these two equalities implies  $\sigma_1(T) = \frac{1}{2}$  and then the second equality is not satisfied.

(e) No mixed strategy equilibrium where 1 uses T and B and 2 uses M and R.

For this we would need:

$$2\sigma_2(M) + 3(1 - \sigma_2(M)) = 7\sigma_2(M) + 4(1 - \sigma_2(M))$$

and

$$7\sigma_1(T) + 2(1 - \sigma_1(T)) = 6\sigma_1(T) + 5(1 - \sigma_1(T))$$

But these equalities imply  $\sigma_1(T) = \frac{3}{4}$  and  $\sigma_2(M) = -\frac{1}{4} < 0$ , which is a contradiction.

- (f) No mixed strategy equilibrium where 1 uses T and B and 2 uses L and M.

For this we would need:

$$7\sigma_2(L) + 2(1 - \sigma_2(L)) = 2\sigma_2(L) + 7(1 - \sigma_2(L))$$

and

$$2\sigma_1(T) + 7(1 - \sigma_1(T)) = 7\sigma_1(T) + 2(1 - \sigma_1(T))$$

But these equalities imply  $\sigma_1(T) = \frac{1}{2}$  and  $\sigma_2(L) = \frac{1}{2}$ . But then the payoff to strategy R is bigger than that for L and M, as

$$6\sigma_1(T) + 5(1 - \sigma_1(T)) = \frac{11}{2} > 7\sigma_1(T) + 2(1 - \sigma_1(T)) = \frac{9}{2},$$

which is a contradiction.

(g) There is a mixed strategy equilibrium where 1 uses T and B and 2 uses L and R.

For this we need:

$$7\sigma_2(L) + 3(1 - \sigma_2(L)) = 2\sigma_2(L) + 4(1 - \sigma_2(L))$$

and

$$2\sigma_1(T) + 7(1 - \sigma_1(T)) = 6\sigma_1(T) + 5(1 - \sigma_1(T))$$

These equalities imply  $\sigma_1(T) = \frac{1}{3}$  and  $\sigma_2(L) = \frac{1}{6}$ . In this case the payoff to strategy M is lower than that for L and R, as

$$6\sigma_1(T) + 5(1 - \sigma_1(T)) = \frac{16}{3} > 7\sigma_1(T) + 2(1 - \sigma_1(T)) = \frac{11}{3}.$$

## Alternative definition of Nash equilibrium

Let

$$B_i(\sigma_{-i}) = \left\{ \sigma_i \in \Sigma_i \mid u_i(\sigma_i, \sigma_{-i}) \geq u_i(\sigma'_i, \sigma_{-i}) \quad \forall \sigma'_i \in \Sigma_i \right\}$$

Then, it is easy to see  $\sigma^*$  is a Nash equilibrium if

$$\sigma_i^* \in B_i(\sigma_{-i}^*) \quad \forall i \in P$$

Also, define  $B(\sigma) = (B_1(\sigma_{-1}), \dots, B_I(\sigma_{-I}))$ . Then  $\sigma^*$  is a Nash equilibrium if

$$\sigma^* \in B(\sigma^*)$$

That is, a Nash equilibrium is a fixed point of  $B(\cdot)$ .



**Theorem 2 (Kakutani)**  $B : \Sigma \rightarrow \Sigma$  has a fixed point if:

1.  $\Sigma$  is a compact, convex, nonempty subset of a Euclidean space.
2.  $B(\sigma)$  is nonempty for all  $\sigma$ .
3.  $B(\sigma)$  is convex for all  $\sigma$ .
4.  $B(\cdot)$  is upper hemi-continuous (alternatively, let any sequence in the domain  $\sigma^n \rightarrow \sigma$ , and any sequence in the range  $\hat{\sigma}^n \rightarrow \hat{\sigma}$  with  $\hat{\sigma}^n \in B(\sigma^n)$ , then if  $\hat{\sigma} \in B(\sigma)$ ,  $B(\cdot)$  is upper-hemicontinuous) .

## Nash equilibrium: existence (3/6)



**Corollary 3** *All finite games have a Nash equilibrium.*

**Proof.** *All we have to show is that conditions 1,2,3 and 4 of previous theorem hold.*

1.  $\Sigma$  obviously nonempty, and is closed and bounded, thus compact.
2.  $u_i(., \sigma_{-i})$  is a continuous function (linear). By Weierstrass theorem a continuous function in a compact set always has a maximum.
3. Suppose  $\sigma' \in B(\sigma)$  and  $\sigma'' \in B(\sigma)$ . Then we must have that

$$\begin{aligned}u_i(\sigma'_i, \sigma_{-i}) &\geq u_i(\sigma_i, \sigma_{-i}) \quad \forall \sigma_i \in \Sigma_i \\u_i(\sigma''_i, \sigma_{-i}) &\geq u_i(\sigma_i, \sigma_{-i}) \quad \forall \sigma_i \in \Sigma_i\end{aligned}$$

thus

$$\lambda u_i(\sigma'_i, \sigma_{-i}) + (1-\lambda)u_i(\sigma''_i, \sigma_{-i}) = u_i(\lambda\sigma'_i + (1-\lambda)\sigma''_i, \sigma_{-i}) \geq u_i(\sigma_i, \sigma_{-i}) \quad \forall \sigma_i \in \Sigma_i$$

4. Suppose not, then  $\exists(\hat{\sigma}^n, \sigma^n) \rightarrow (\hat{\sigma}, \sigma)$  with  $\hat{\sigma}^n \in B(\sigma^n)$  but  $\hat{\sigma} \notin B(\sigma)$ . Thus there must be some  $i \in P$  with  $\hat{\sigma}_i \notin B_i(\sigma_{-i})$ . Thus, there is some  $\varepsilon > 0$  and some  $\sigma'_i$  with  $u_i(\sigma'_i, \sigma_{-i}) \geq u_i(\hat{\sigma}_i, \sigma_{-i}) + 3\varepsilon$  (a). Also, by continuity of  $u_i(\cdot)$  and since  $(\hat{\sigma}^n, \sigma^n) \rightarrow (\hat{\sigma}, \sigma)$  we must have that there is  $n$  large enough that:

$$u_i(\sigma'_i, \sigma_{-i}^n) > u_i(\sigma'_i, \sigma_{-i}) - \varepsilon$$

Now by (a) we must have

$$u_i(\sigma'_i, \sigma_{-i}) - \varepsilon > u_i(\hat{\sigma}_i, \sigma_{-i}) + 2\varepsilon$$

and continuity again

$$u_i(\hat{\sigma}_i, \sigma_{-i}) + 2\varepsilon > u_i(\hat{\sigma}_i^n, \sigma_{-i}^n) + \varepsilon$$

which contradicts  $\hat{\sigma}_i^n \in B(\sigma_{-i}^n)$  ■

**Corollary 4** *All infinite games have a Nash equilibrium provided that.*

(a)  $S_i$  are nonempty compact, convex subsets of a Euclidean space.

(b)  $u_i(\cdot)$  is continuous in  $S$  and quasi-concave in  $s_i$

**Theorem 5 Proof.** 1. True by (a).

2.  $u_i(\cdot)$ ,  $S$  is compact by (a). By Weierstrass theorem a continuous function in a compact set always has a maximum.

3. By definition of quasi-concavity of  $B(\cdot)$  we have that for any  $s'_i$  and  $s''_i$  with:

$$\begin{aligned}u_i(s'_i, s_{-i}) &\geq u_i(s_i, s_{-i}) \quad \forall s_i \in S_i \\u_i(s''_i, s_{-i}) &\geq u_i(s_i, s_{-i}) \quad \forall s_i \in S_i\end{aligned}$$

we must have that:

$$u_i(\lambda s'_i + (1 - \lambda)s''_i, s_{-i}) \geq u_i(s_i, s_{-i}) \quad \forall s_i \in \Sigma_i$$

so  $B(s)$  is convex for all  $s$ .

4.  $u_i(\cdot)$  is continuous by (b). ■

**Remark 6** *When  $u_i$  is continuous but not quasi-concave, mixed strategies can give an equilibrium.*

*The proof needs more machinery but is very similar.*

*$S_i$  need not be convex now, as mixed strategies convexify strategy set.*

*Also mixed strategies make payoff linear and continuous, and best responses convex-valued.*

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## Chapter 1

### Games in Strategic Form - Nash equilibrium

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