

# Repeated games

## 1: Finite repetition

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# Finitely repeated games

- A **finitely** repeated game is a dynamic game in which a simultaneous game (the **stage game**) is played finitely many times, and the result of each stage is **observed** before the next one is played.
- Example: Play the prisoners' dilemma several times. The stage game is the simultaneous prisoners' dilemma game.

# Results

- If the stage game (the simultaneous game) has **only one NE** the repeated game has **only one SPNE**: In the SPNE players' play the strategies in the NE in each stage.
- If the stage game has **2 or more NE**, one can find a SPNE where, at some stage, players play a strategy that is **not part of a NE** of the stage game.

# The prisoners' dilemma repeated twice

- Two players play the same **simultaneous** game twice, at  $t = 1$  and at  $t = 2$ .
- After the first time the game is played (after  $t = 1$ ) the result is **observed** before playing the second time.
- The payoff in the repeated game is the **sum** of the payoffs in each stage ( $t = 1, t = 2$ )
- Which is the SPNE?

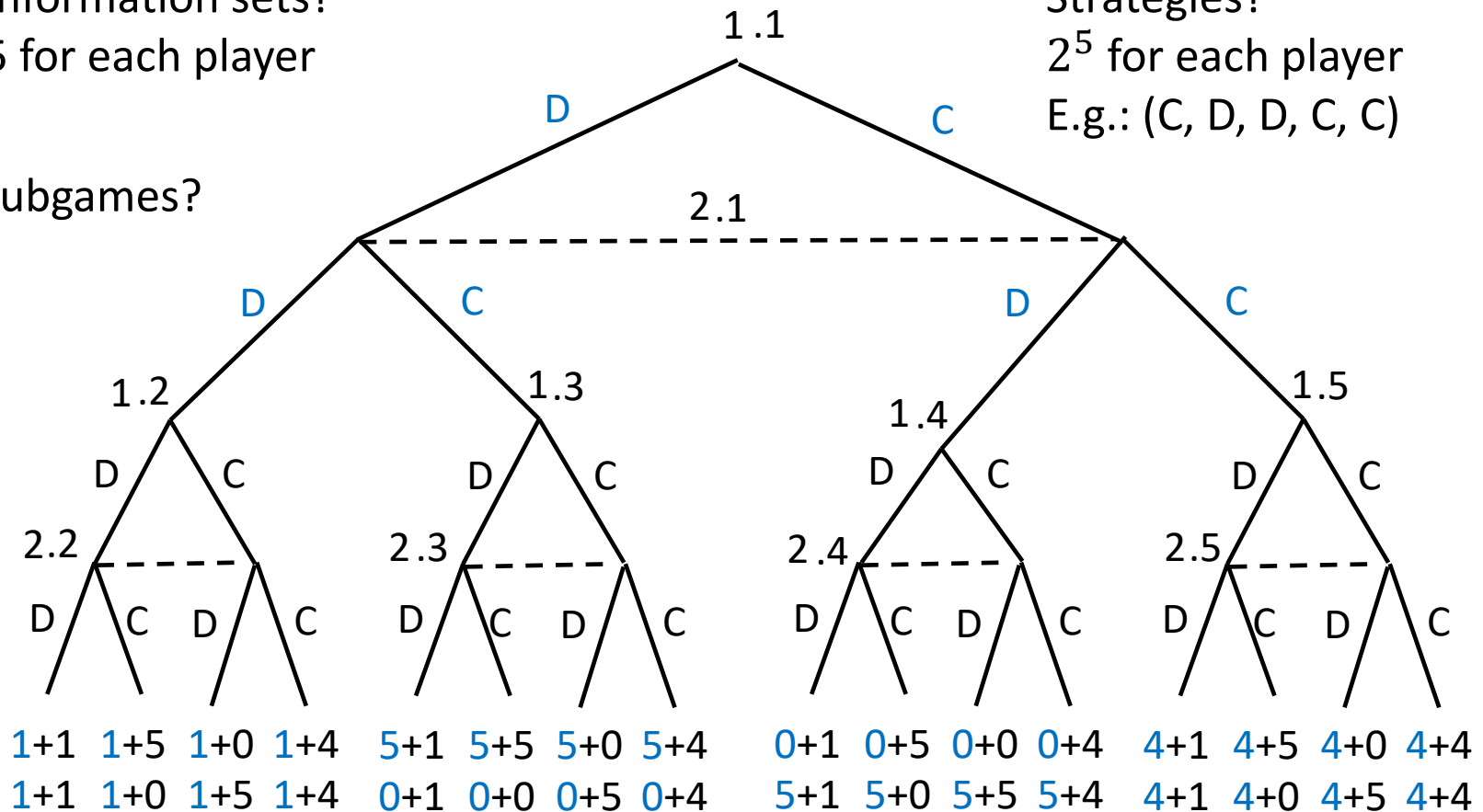
		Player 2	
		D	C
Player 1	D	1, 1	5, 0
	C	0, 5	4, 4

# The prisoners' dilemma repeated twice

Information sets?  
5 for each player

Strategies?  
 $2^5$  for each player  
E.g.: (C, D, D, C, C)

Subgames?  
5

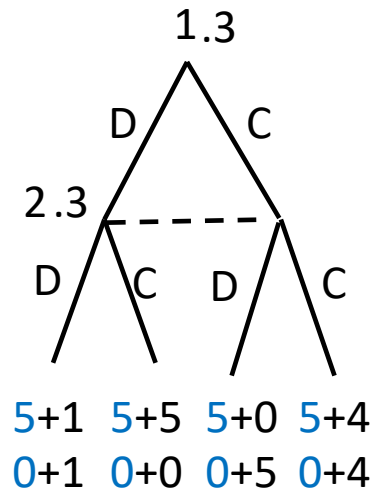


# The prisoners' dilemma repeated twice

Let's find the NE in the subgames.

Consider the subgame starting at 1.3.

Other subgames are solved similarly.



Player 1

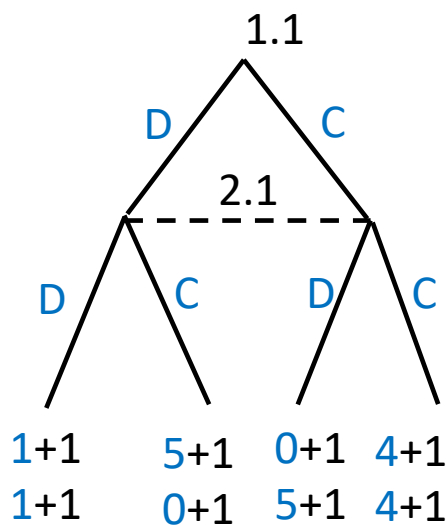
		Player 2	
		D	C
Player 1	D	<u>5+1</u> , <u>0+1</u>	5+5, 0+0
	C	5+0, <u>0+5</u>	5+4, 0+4

		Player 2	
		D	C
Player 1	D	<u>1</u> , <u>1</u>	<u>5</u> , 0
	C	0, <u>5</u>	4, 4

Adding or not the **payoffs of the first stage** does not change the game. In each subgame we have the **same** prisoners' dilemma.

# The prisoners' dilemma repeated twice

- In each of the four subgames there is **only one NE**: (D, D).
- Using backward induction, we **substitute** the payoffs in the equilibria ((1, 1) in all of them) and proceed to solve the first stage.



Player 1

		Player 2	
		D	C
Player 1	D	<u>1+1</u> , <u>1+1</u>	<u>5+1</u> , 0+1
	C	0+1, <u>5+1</u>	4+1, 4+1

Player 1

		Player 2	
		D	C
Player 1	D	<u>1</u> , <u>1</u>	<u>5</u> , 0
	C	0, <u>5</u>	4, 4

Again, we have a **prisoners' dilemma**, with NE: (D, D).

The **only** SPNE is ((D, D, D, D, D), (D, D, D, D, D)).

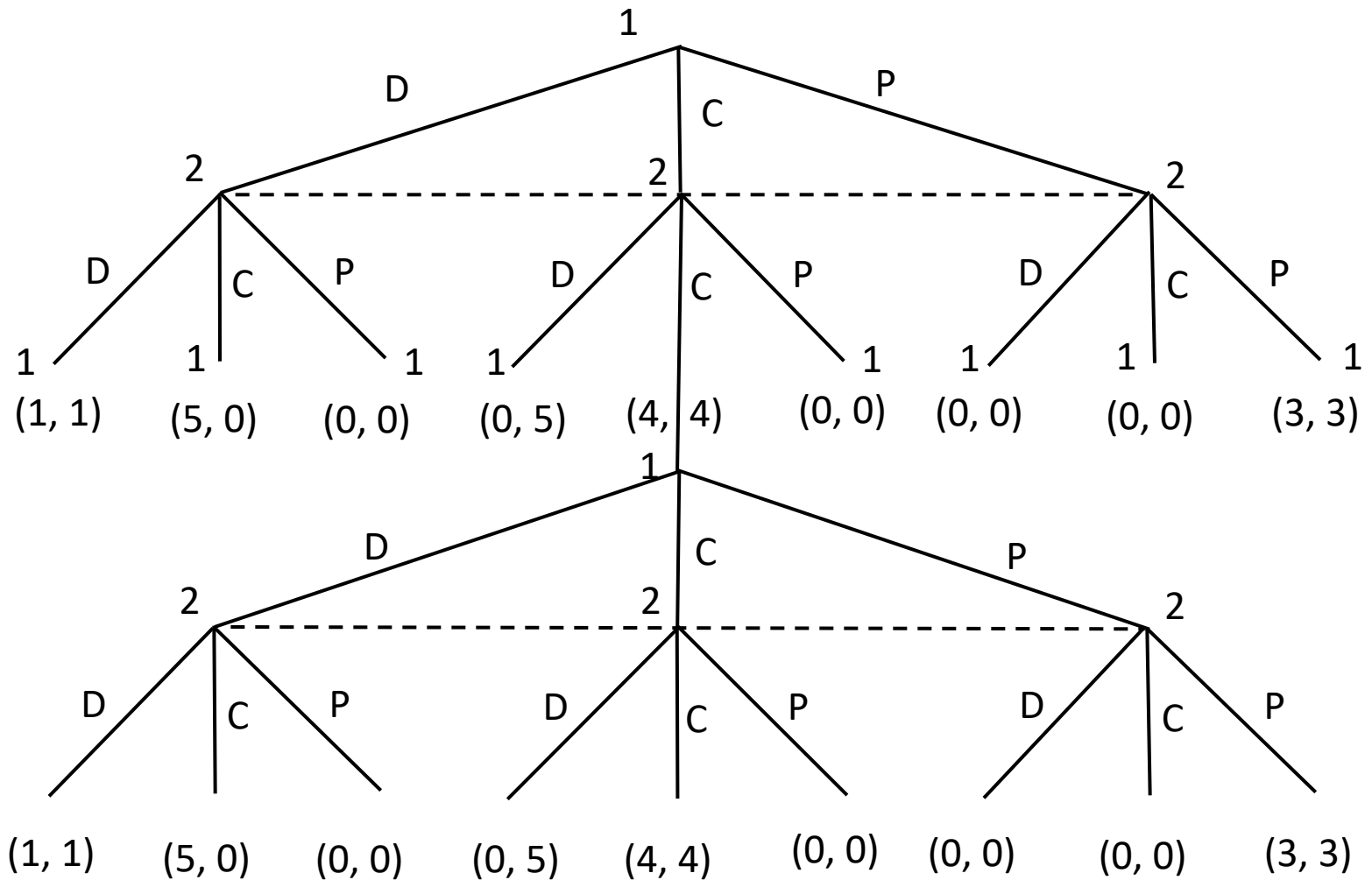
# The augmented prisoners' dilemma

		Player 2		
		D	C	P
Player 1	D	<u>1</u> , <u>1</u>	<u>5</u> , 0	0, 0
	C	0, <u>5</u>	4, 4	0, 0
	P	0, 0	0, 0	<u>3</u> , <u>3</u>

- Let's play **twice** the game before. It is the prisoners' dilemma with an **added** strategy.
- (D, D) is still a Nash equilibrium.
- There is also a new Nash equilibrium: (P, P).
- Could the play of (C, C) be part of a SPNE?
- In the second stage we may not have the same NE in all subgames and, thus, the game in the first stage may not be the same as the stage game.



# The augmented prisoners' dilemma



# The augmented prisoners' dilemma

- To sustain cooperation (C, C) in the first stage, let's **play with the NE in the second one**:
  - If there has been cooperation in the first stage, in the second stage play the **NE with the best** payoffs: (P, P).
  - If there has not been cooperation, in the second stage play the **NE with worse** payoffs: (D, D).
- Formally:
  - At  $t = 1$ : (C, C).
  - At  $t = 2$ : (P, P) if (C, C) was played at  $t = 1$ ,  
(D, D) if (C, C) was not played at  $t = 1$ .

# The augmented prisoners' dilemma

At  $t = 1$ : (C, C).

At  $t = 2$ : (P, P) if (C, C) was played at  $t = 1$ ,  
(D, D) if (C, C) was not played  $t = 1$ .

- Let us check if the above strategy is a SPNE:
  - It prescribes a Nash equilibrium in **each subgame of the second stage**: (P, P) or (D, D).
  - It remains to check whether it is a NE of the **whole game**.
- If the players **follow** the strategy, each one gets  $4 + 3 = 7$ .
- If one player **deviates**:
  - **In the second stage there is no profit** in deviating (recall that players play a NE).
  - **In the first stage, the best deviation from (C, C) is D**. In this first stage, the player that deviates gets 5, but then in the second stage the game reaches a subgame in which the equilibrium is (D, D), with payoff 1.
  - The total for the deviating player is  $5 + 1 = 6$ , less than the 7 in case of no deviation.
- Thus, the strategy is a **SPNE**.

# The augmented prisoners' dilemma

Recall:

AT  $t = 2$ : (P, P) if (C, C) was played at  $t = 1$ ,

(D, D) if (C, C) was not played at  $t = 1$ .

		Player 2		
		D	C	P
Player 1	D	1 + 1, 1 + 1	5 + 1, 0 + 1	0 + 1, 0 + 1
	C	0 + 1, 5 + 1	4 + 3, 4 + 3	0 + 1, 0 + 1
	P	0 + 1, 0 + 1	0 + 1, 0 + 1	3 + 1, 3 + 1

		Player 2		
		D	C	P
Player 1	D	<u>2</u> , <u>2</u>	6, 1	1, 1
	C	1, 6	<u>7</u> , <u>7</u>	1, 1
	P	1, 1	1, 1	<u>4</u> , <u>4</u>

# The augmented prisoners' dilemma

- For any **combination** of NE in the subgames of the second stage, there will be different SPNE.
- How many combinations are there?
- It's 9 subgames, with two NE in pure strategies in each one. Thus, there are  $2^9 = 512$  different combinations.
- It will be very complicated to analyze all of them. Instead, we can look for **interesting equilibria**, like the one we just found.
- Let's see three more **examples** of SPNE:
  - Play (D, D) always (at  $t = 1$  and at  $t = 2$  in all subgames).
  - Play (P, P) always.
  - Play (D, D) at  $t = 1$  and (P, P) at  $t = 2$  in all subgames.
- There are **many more**. In fact, in these games we typically study the **set of payoffs** that can be obtained in a SPNE.

# The chicken game

	K	S
K	0, 0	5, 1
S	1, 5	4, 4

- The game has **3 NE**:  $\left\{ (S, K), (K, S), \left( \frac{1}{2} [K] + \frac{1}{2} [S], \frac{1}{2} [K] + \frac{1}{2} [S] \right) \right\}$ , with respective payoffs: (1, 5), (5, 1), (2,5, 2,5).
- Repeat the game **twice**. Can we sustain **(S, S) in the first stage**?
- Define an appropriate strategy:
  - At  $t = 1$ : (S, S).
  - At  $t = 2$ :  $\left( \frac{1}{2} [K] + \frac{1}{2} [S], \frac{1}{2} [K] + \frac{1}{2} [S] \right)$  if either (S, S) or (K, K) was played at  $t = 1$ ,  
 (S, K) if (K, S) was played at  $t = 1$ ,  
 (K, S) if (S, K) was played at  $t = 1$ .
- Observe that this strategy **only punishes the unilateral deviator**.
- Utility if there are **no deviations**:  $4 + 2,5 = 6,5$ .
- Utility for  $i$  if he **deviates at  $t = 1$** :  $5 + 1 = 6$ .
- The trigger strategy is a SPNE.