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Foundations for contest success functions

**Authors:** Luis Corchón · Matthias Dahm

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# Foundations for contest success functions

Luis Corchón · Matthias Dahm

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**Abstract** In the literature, the outcome of contests is either interpreted as win probabilities or as shares of the prize. With this in mind, we examine two approaches to contest success functions (CSFs). In the first, we analyze the implications of contestants' incomplete information concerning the 'type' of the contest administrator. While in the case of two contestants this approach can rationalize prominent CSFs, we show that it runs into difficulties when there are more agents. Our second approach

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7 interprets CSFs as sharing rules and establishes a connection to bargaining and claims  
 8 problems which is independent of the number of contestants. Both approaches provide  
 9 foundations for popular CSFs and guidelines for the definition of new ones.

10 **Keywords** Endogenous contests · Contest success function ·  
 11 Nash bargaining solution · Bargaining with claims

12 **JEL Classification** C72 · D72 · D74

## 13 1 Introduction

14 “The strategic approach also seeks to combine axiomatic cooperative solutions and  
 15 non-cooperative solutions. Roger Myerson recently named this task the ‘Nash pro-  
 16 gram’.” (Rubinstein 1985, p. 1151)

17 A contest is a game in which players exert effort to win a certain prize. Contests  
 18 have been used to analyze a variety of situations including lobbying, rent-seeking and  
 19 rent-defending contests, advertising, litigation, political campaigns, conflict, patent  
 20 races, arms races, sports events or R&D competition. A crucial determinant for the  
 21 equilibrium predictions of contests is the specification of the so-called contest success  
 22 function (CSF) which relates the players’ efforts and win probabilities. Justifications  
 23 for a particular CSF can be twofold. A justification can be on normative grounds,  
 24 because it is the unique CSF fulfilling certain axioms, or essential properties. A jus-  
 25 tification can also be positive when it can be shown that the CSF arises from the  
 26 strategic interaction of players, thereby yielding a description of situations when it  
 27 can be expected to be realistic. The purpose of the present paper is to contribute to our  
 28 understanding of CSFs in both dimensions.

29 Formally, a CSF associates, to each vector of efforts  $\mathbf{G}$ , a lottery specifying for each  
 30 agent a probability  $p_i$  of getting the object. That is,  $p_i = p_i(\mathbf{G})$  is such that, for each  
 31 contestant  $i \in N := \{1, \dots, n\}$ ,  $p_i(\mathbf{G}) \geq 0$ , and  $\sum_{i=1}^n p_i(\mathbf{G}) = 1$ .

32 The canonical example of a contest situation is rent-seeking. In a pioneering paper,  
 33 Tullock (1980) proposed a special form of the CSF, namely, given a positive scalar  
 34  $R$ ,

$$35 \quad p_i = \frac{G_i^R}{\sum_{j=1}^n G_j^R}, \quad \text{for } i = 1, \dots, n. \quad (1)$$

36 Gradstein (1995, 1998) postulated the following variation of this form, where, given  
 37  $q_i > 0$  for all  $i \in N$ ,

$$38 \quad p_i = \frac{G_i q_i}{\sum_{j=1}^n G_j q_j}, \quad \text{for } i = 1, \dots, n. \quad (2)$$

39 A generalization that comprises both previous functional forms is, given  $a_i \geq 0$  for  
 40 all  $i \in N$ ,

$$41 \quad p_i = \frac{G_i^R q_i + a_i}{\sum_{j=1}^n (G_j^R q_j + a_j)}, \quad \text{for } i = 1, \dots, n. \quad (3)$$

42 A different functional form, the logit model, was proposed by [Hirshleifer \(1989\)](#),  
 43 where, given a positive scalar  $k$ ,

$$44 \quad p_i = \frac{e^{kG_i}}{\sum_{j=1}^n e^{kG_j}}, \quad \text{for } i = 1, \dots, n. \quad (4)$$

45 Note that the four expressions (1)–(4) are specific instances of the following functional  
 46 form

$$47 \quad p_i = \frac{f_i(G_i)}{\sum_{j=1}^n f_j(G_j)}, \quad \text{for } i = 1, \dots, n. \quad (5)$$

48 The so-called effectivity functions  $f_i$  are usually interpreted as determining how  
 49 ‘effective’ agent  $i$ ’s effort is in affecting the win probability of agent  $i$ . Most papers  
 50 dealing with contest models in the literature analyze a CSF, which is a special case  
 51 of the additive form in (5) ([Nitzan 1994](#); [Konrad 2007](#)). Consequently, the present  
 52 paper will be mainly concerned with deriving foundations for CSFs of this form.  
 53 Notice, for later reference, that in (5) the win probability of any contestant is respon-  
 54 sive to changes in the efforts of all other contestants, if the  $f_i$  are strictly increas-  
 55 ing.

56 However, there are also some CSFs in the literature, which are not special cases  
 57 of the form in (5). The first two consider the case of two contestants and build on the  
 58 idea that only differences in effort should matter—an idea introduced by [Hirshleifer](#)  
 59 in (4). [Baik \(1998\)](#) proposed the following form, given a positive scalar  $\sigma$ ,

$$60 \quad p_1 = \frac{\sigma G_1 - G_2}{\sigma G_1 - G_2 + 1} \quad \text{and} \quad p_2 = 1 - p_1. \quad (6)$$

61 [Che and Gale \(2000\)](#) postulate the following piecewise linear difference-form

$$62 \quad p_1 = \max \left\{ \min \left\{ \frac{1}{2} + \sigma(G_1 - G_2), 1 \right\}, 0 \right\} \quad \text{and} \quad p_2 = 1 - p_1. \quad (7)$$

63 Recently, [Alcalde and Dahm \(2007\)](#) proposed a CSF in which relative differences  
 64 matter. Given an ordered vector of efforts such that  $G_1 \geq G_2 \geq \dots \geq G_n$  and a  
 65 positive scalar  $R$ , the serial CSF is defined as

$$66 \quad p_i = \sum_{j=i}^n \frac{G_j^R - G_{j+1}^R}{j \cdot G_1^R}, \quad \text{for } i = 1, \dots, n \quad \text{with} \quad G_{n+1} = 0. \quad (8)$$



In the literature, the outcome of contests has been interpreted to capture two different situations: as win probabilities or as shares of the prize.<sup>1</sup> With this in mind, we examine two approaches to CSFs.

In the first, we postulate the existence of a contest administrator who allocates the prize to one of the contestants. However, contestants have incomplete information about the type of the contest administrator. We show that this approach can generate CSFs for any number of contestants. However, while in the case of two contestants this approach can rationalize a large class of CSFs, we show that it runs into difficulties when there are more agents.

Our second approach interprets CSFs as sharing rules, and establishes a connection to bargaining and claims problems, which is independent of the number of contestants. The analysis exploits the observation that these problems are mathematically related—but not equivalent—to the problem of assigning win probabilities in contests. A main result here follows [Dagan and Volij \(1993\)](#), and shows that the class of CSFs given in (5) can be understood as the weighted Nash bargaining solution where efforts represent the weights of the agents. We turn then to the framework of bargaining with claims ([Chun and Thomson 1992](#)) to incorporate explicitly the contestants' efforts in the description of the problem. This allows to associate prominent solution concepts in this framework to the previously mentioned class of CSFs and to a generalized version of Che and Gale's difference-form contest (7).

Both approaches provide foundations for popular CSFs and guidelines for the definition of new ones. In our view, both types of foundations complement each other nicely. For instance, we show that (7) can be understood, on one hand, as contestants trying to sway away the contest administrator's decision in a setting analogous to the model of a circular city by [Salop \(1979\)](#). On the other, we show that this CSF is also related to the claim-egalitarian solution ([Bossert 1993](#)). Both approaches lend support to an extension of this CSF to three contestants of the following form. Let  $G_1 \geq G_2 \geq G_3$  and  $a$  and  $b$  be positive scalars. If  $G_1 - G_3 \geq a$  then  $p_3 = 0$  and the other contestants obtain win probabilities as in (7). Otherwise, let

$$p_i = \frac{1}{3} + b(2G_i - G_j - G_k), \quad \text{for } i = 1, 2, 3 \quad \text{and } i \neq j, k. \quad (9)$$

However, the requirement that for  $n = 2$  the CSF reduces to (7) implies that  $(a, b) = ((3\sigma)^{-1}, \sigma/2)$  in the first and  $(a, b) = ((2\sigma)^{-1}, 2\sigma/3)$  in the second approach. This underlines that the appropriate extension depends on the application and institutional details the contest model is intended to capture.

Foundations for CSFs have been reviewed by [Garfinkel and Skaperdas \(2007\)](#) and [Konrad \(2007\)](#). The most systematic approach has been normative and the seminal paper is that by [Skaperdas \(1996\)](#). He proposed five axioms and showed that they are equivalent to assuming a CSF of the form given in (5) with  $f_i(\cdot) = f(\cdot)$  for all  $i \in N$ , where  $f(\cdot)$  is a positive increasing function of its argument. Skaperdas also showed that if in addition to the other five axioms the CSF is assumed to be homogeneous

<sup>1</sup> A prominent example for the latter is [Wärneryd \(1998\)](#). He analyzes a contest among jurisdictions for shares of the GNP and compares different types of jurisdictional organization.

of degree zero in  $G$  then we obtain (1).<sup>2</sup> Our paper contributes to this literature indirectly by making connections to related problems, which are well understood from a normative point of view. For instance, we establish a relationship between Che and Gale's difference-form CSF (7) and the principle of equal sacrifice.

As for the positive approach, we are not aware of any work understanding CSFs as sharing rules as our second approach does.<sup>3</sup> However, our first approach is related to other works. Assume that efforts are a noisy predictor of performance in the contest. When noise enters additively in performance and is distributed as the extreme value distribution, we obtain the logit specification (McFadden (1974)). This procedure was generalized by Lazear and Rosen (1981) and Dixit (1987) to general distributions.<sup>4</sup> Our approach that differs from these papers by changing performance to the broader concept of utility and using a uniformly distributed one-dimensional random variable.

Epstein and Nitzan (2006) partially rationalize CSFs by analyzing how a contest administrator rationally decides whether to have a contest and if a contest takes place how he chooses among a fixed set of CSFs. In contrast, in our approach, the administrator chooses deterministically, but the contestants face a CSF because of their uncertainty about the type of the administrator.

## 2 External decider

### 2.1 Two contestants

Assume that one person has to decide to award a prize to one of two contestants. In this situation we have in mind that the contestants are uncertain about a characteristic of the decider that is relevant for his decision. So contestants exert effort without knowing the realization of the characteristic and then the decision-maker decides whom to give the prize based both on the contestants' efforts and his type.

Let  $\Theta$  be the set of states of the world. Let  $\theta$  be an arbitrary element of  $\Theta$ . We assume that  $\Theta = [0, 1]$  and that  $\theta$  is uniformly distributed. Let  $V_i$  be the decider's payoff if the prize is awarded to contestant  $i = 1, 2$ .  $V_i$  is assumed to depend on the state of the world, i.e.  $V_i = V_i(\theta)$ . This may reflect the uncertainty in the contestants' minds

<sup>2</sup> An extension of Skaperdas' result to nonanonymous CSFs is given by Clark and Riis (1998). Skaperdas also axiomatized the logit model (4).

<sup>3</sup> Anbarci et al. (2002) present a model in which a two-party conflict over a resource can either be settled through bargaining over the resource or through a contest. The contest defines the disagreement point of the bargaining problem to which three different bargaining solutions are applied (see also Esteban and Sákovicš 2006). In contrast, in our framework, we interpret bargaining to be over win probabilities and derive CSFs as bargaining rules.

<sup>4</sup> Hillman and Riley (1989) came close to the idea of a contest administrator. They propose a 'political impact' function that reflects the influence of a player as a function of her effort and a random variable. They notice that for two agents it is possible to specify a functional form for this function, which yields the Tullock probability function (see also Hirschleifer and Riley 1992). This was generalized by Jia (2007) to  $n > 2$ . In related work, Fullerton and McAfee (1999) and Baye and Hoppe (2003) offer microfoundations for a subset of CSFs of the form in (1) in the context of innovation tournaments and patent races following an analogous procedure.

135 about the preferences of the decider. We will assume the following single-crossing  
136 property.

137 (SC)  $V_1(\theta)$  is decreasing in  $\theta$  and  $V_2(\theta)$  is strictly increasing in  $\theta$ .

138 Taking into account the efforts, let  $U_i(V_i(\theta), G_i)$  be the decider's payoff if the prize  
139 is awarded to contestant  $i = 1, 2$ . This function is assumed to be increasing in both  
140 arguments and for simplicity we will write  $U_i(\theta, G_i)$ . For the sake of interpretation,  
141 let  $G_i$  be interpreted as the level of advertizement (resp. quality) made (resp. provided)  
142 by Contestant  $i = 1, 2$ . Let

$$143 \quad \theta' = \begin{cases} 1 & \text{if } U_1(\theta, G_1) > U_2(\theta, G_2), \quad \forall \theta \in \Theta \\ 0 & \text{if } U_1(\theta, G_1) < U_2(\theta, G_2), \quad \forall \theta \in \Theta \\ \{\theta | U_1(\theta, G_1) = U_2(\theta, G_2)\} & \text{otherwise.} \end{cases} \quad (10)$$

144 Under our assumptions  $\theta'$  is well-defined and unique. Moreover,  $\theta'$  equals  $p_1$ , the  
145 probability that Contestant 1 gets the prize. We now provide several examples in  
146 which we solve for  $p_1$  as a function of  $G_1$  and  $G_2$ . This way we obtain the CSF as  
147 arising from the maximization of the payoff function of the decider.

148 In these examples,  $V_i(\theta)$  enters either additively [in the spirit of [McFadden \(1974\)](#)]  
149 or multiplicatively [as in [Hillman and Riley \(1989\)](#)]. In Examples 1 and 2, the effect  
150 of a contestant's advertizement is completely separated from the decider's bias. The  
151 function  $U_i(\theta, G_i)$  is additively separable in both arguments. Here, the merit of an  
152 alternative in the decider's eyes might be positive even when advertizing is zero,  
153 and vice versa. Moreover, the marginal product of advertizing is independent of the  
154 decider's bias. This contrasts with the multiplicative form of Example 3 in which (i)  
155 a prerequisite for the merit of an alternative is both that the decider likes it (at least a  
156 little) and that advertizing is positive; and (ii) an increase of the decider's bias raises  
157 the marginal product of advertizing. Example 4 is a combination of these two extreme  
158 cases in the sense that for one contestant the relationship is multiplicative, while for  
159 the other the effect of advertizing is independent of the bias.

160 *Example 1* Let  $U_1(\theta, G_1) = V_1(\theta) + a_1 G_1$  and  $U_2(\theta, G_2) = V_2(\theta) + a_2 G_2$ , where  
161  $a_1, a_2 > 0$ . Thus,  $a_1 G_1 - a_2 G_2 = V_2(\theta) - V_1(\theta) \equiv z(\theta)$ , say. Since  $z(\cdot)$  is invert-  
162 ible, we get  $p_1 = z^{-1}(a_1 G_1 - a_2 G_2)$ , which is the form in (6) considered by [Baik](#)  
163 (1998).<sup>5</sup> Notice that this procedure is identical to the one used in models of spatial  
164 differentiation to obtain the demand function (see [Hotelling 1929](#)).

165 *Example 2* Let  $U_1(\theta, G_1) = \theta + 2\sigma G_1 - 1/2$  and  $U_2(\theta, G_2) = -\theta + 2\sigma G_2 + 1/2$ ,  
166 where  $\sigma$  is a positive scalar. In this case, it is easily calculated that

$$167 \quad p_1 = \max \{ \min \{ 1/2 + \sigma(G_1 - G_2), 1 \}, 0 \}.$$

168 We obtain (7), the family of difference-form CSFs analyzed by [Che and Gale \(2000\)](#).

<sup>5</sup> Alternatively, we may assume that the payoff function of the decider is  $U_i = V_i(\theta) - a_j G_j$ ,  $i \neq j$ , reflecting the disutility received from the effort made by Contestant 2, if the prize is awarded to Contestant 1. The same applies to Example 2 and 3 by taking  $U_1 = (1 - \theta)/f_2(G_2)$  and  $U_2 = \theta/f_1(G_1)$ .

169 *Example 3* Let  $U_1(\theta, G_1) = (1 - \theta)f_1(G_1)$  and  $U_2(\theta, G_2) = \theta f_2(G_2)$ . Here, we  
 170 obtain  $p_1 = f_1(G_1)/(f_1(G_1) + f_2(G_2))$ . This is the additive CSF (5) for  $n = 2$ .

171 *Example 4* Let  $U_1(\theta, G_1) = f_1(G_1)$  and  $U_2(\theta, G_2) = 2\theta f_2(G_2)$  if  $\theta \leq 1/2$ , whereas  
 172  $U_2(\theta, G_2) = f_2(G_2)/(2(1-\theta))$  if  $1/2 \leq \theta < 1$ . By analogous reasoning as before, we  
 173 obtain  $p_1 = f_1(G_1)/(2f_2(G_2))$  if  $f_1(G_1) \leq f_2(G_2)$  and  $p_1 = 1 - f_2(G_2)/(2f_1(G_1))$   
 174 otherwise. This expression is a generalization of the family of serial contests (8) ana-  
 175 lyzed by Alcalde and Dahm (2007).

176 To derive a general result concerning what kind of CSFs can be derived from the  
 177 maximization of the payoffs of the decider we will now consider the class of CSF,  
 178 which are  $\mathbb{C}^1$  in  $\mathbb{R}_{++}^n$ . This leaves outside our study CSFs like (7), but includes (8)  
 179 when  $n = 2$ .

180 A difficulty in our study is that many well-known CSFs fail to be continuous when  
 181  $G_i = 0$  all  $i$  and constant in its own effort when  $G_j = 0$  all  $j \neq i$ , e.g. (1). A way to  
 182 solve these problems is to stay away from the troublesome boundaries of  $\mathbb{R}_{++}^n$  as we  
 183 do in Definitions 1 and 2.

184 **Definition 1**  $p_i = p_i(\mathbf{G})$  is regular if for all  $\mathbf{G} \in \mathbb{R}_{++}^n$ ,  $\partial p_i(\mathbf{G})/\partial G_i > 0$  and  
 185  $\partial p_i(\mathbf{G})/\partial G_j < 0$  for all  $j \neq i$ .

186 Notice that the CSFs in (1)–(4) and (6) are regular. The one in (5) is regular if we  
 187 assume, as in Szidarovsky and Okuguchi (1997), that  $f'_i(G_i) > 0$  and  $f_i(0) = 0$  for  
 188 all  $i \in N$ . The CSF given in (8) is regular if  $n = 2$ .

189 **Definition 2** The CSF  $\{p_1(\mathbf{G}), p_2(\mathbf{G}), \dots, p_n(\mathbf{G})\}$  is rationalizable if there is a list  
 190 of payoff functions  $U_i(\theta, G_i)$  strictly increasing on  $G_i$ ,  $i = 1, 2, \dots, n$  such that for  
 191 any  $\hat{\mathbf{G}} \in \mathbb{R}_{++}^n$ ,

192 
$$p_i(\hat{\mathbf{G}}) = \text{probability}\{U_i(\theta, \hat{G}_i) > U_j(\theta, \hat{G}_j), \forall j \neq i\}, \quad \text{for } i = 1, \dots, n.$$

193 We need the following assumption:

194 **Assumption 1**  $p_i \rightarrow 1$  when  $G_i \rightarrow \infty$  and  $p_i \rightarrow 0$  when  $G_i \rightarrow 0$ .

195 It is easy to check that Tullock’s CSF (1) satisfies Assumption 1 (A.1 in the sequel).  
 196 Also the additive CSF (5) satisfies A.1 when  $f_i(G_i)$  are strictly positive for strictly  
 197 positive values of efforts,  $f_i \rightarrow \infty$  when  $G_i \rightarrow \infty$  and  $f_i \rightarrow 0$  when  $G_i \rightarrow 0$ . It  
 198 is fulfilled by the serial CSF in (8) and the form in (6) includes cases where A.1 is  
 199 satisfied. Now we can prove the following:

200 **Proposition 1** If A.1 holds and  $p_1(G_1, G_2)$  is regular, it is rationalizable by a pair  
 201 of payoff functions fulfilling the single crossing condition. If  $p_1(G_1, G_2)$  is ratio-  
 202 nalizable by a pair of payoff functions fulfilling the single crossing condition and  
 203  $\partial p_i(\mathbf{G})/\partial G_j \neq 0$  for all  $i, j$ , it is regular.

204 *Proof* Suppose  $p_1(G_1, G_2)$  is regular. Notice that this implies that for any  $\mathbf{G} \in \mathbb{R}_{++}^2$ ,  
 205  $p_i \in (0, 1)$ . Let  $f(p_1, G_1, G_2) \equiv p_1 - p_1(G_1, G_2)$ . Fix  $p_1$  and  $G_2$ , say  $\bar{p}_1$  and  $\bar{G}_2$ . By

206 A.1, we have that  $f(\bar{p}_1, G_1, \bar{G}_2) < 0$  for  $G_1$  sufficiently large and  $f(\bar{p}_1, G_1, \bar{G}_2) > 0$   
 207 for  $G_1$  sufficiently close to zero. By the intermediate value theorem, there is a  $G_1$  such  
 208 that  $f(\bar{p}_1, G_1, \bar{G}_2) = 0$ . By the definition of a regular CSF this value of  $G_1$ , say  $\hat{G}_1$ ,  
 209 is unique. This means that there is a unique function  $H$  such that  $G_1 = H(p_1, G_2)$ .  
 210 Since  $\partial f(p_1, G_1, G_2)/\partial G_1 < 0$ , by the implicit function theorem  $H$  is continuous  
 211 in the neighborhood of  $(\bar{p}_1, \bar{G}_2)$ . Since this point is arbitrary,  $H$  is continuous for all  
 212  $(p_1, G_2)$ . Let  $U_1 = G_1$  and  $U_2 = H(\theta, G_2)$ . Because  $p_1(G_1, G_2)$  is regular,  $H$  is  
 213 strictly increasing on  $\theta$  and  $G_2$ . Also  $U_1$  is strictly increasing on  $G_1$  and constant on  
 214  $\theta$  so, the SC assumption holds. By construction,  $\theta'$  (as defined in Eq. 10) equals  $p_1$ ;  
 215 thus,  $p_1(G_1, G_2)$  is rationalizable.

216 Assume now that  $p_1(G_1, G_2)$  is rationalizable by a list of payoff functions fulfilling  
 217 the single crossing condition (SC). Rationalizability implies that for any  $(\hat{G}_1, \hat{G}_2)$  we  
 218 have  $p_1(\hat{G}_1, \hat{G}_2) = \theta'$  (as defined in Eq. 10). Moreover, as  $U_1$  is strictly increasing on  
 219  $G_1$  and by the single crossing condition (SC)  $U_2$  is strictly increasing on  $\theta$ , we have  
 220 that  $p_1$  is strictly increasing in  $G_1$ . The opposite holds when  $G_2$  is increased; so, the  
 221 result follows from  $\partial p_i(\mathbf{G})/\partial G_j \neq 0$ .  $\square$

222 We show now that the condition that the partial derivatives do not vanish cannot be  
 223 dispensed with.

224 *Example 5* Consider the following smooth difference-form contest between two con-  
 225 testants:

$$226 \quad p_1 = \begin{cases} 1 & \text{if } G_1 - G_2 \geq 1 \\ \frac{1}{2} + \frac{1}{2}e^{\left\{ \frac{-(G_1 - G_2 - 1)^2}{1 - (G_1 - G_2 - 1)^2} \right\}} & \text{if } 1 > G_1 - G_2 \geq 0 \\ \frac{1}{2}e^{\left\{ \frac{-(G_1 - G_2)^2}{1 - (G_1 - G_2)^2} \right\}} & \text{if } 0 \geq G_1 - G_2 > -1 \\ 0 & \text{if } -1 \geq G_1 - G_2 \end{cases} \quad \text{and } p_2 = 1 - p_1. \quad (11)$$

227 As in (7), the win probability might be zero—even for positive effort. Contrary to  
 228 (7) it is  $\mathbb{C}^1$ . Notice that for  $|G_1 - G_2| \leq 1$ ,  $p_1$  is strictly monotonic. However,  
 229 when  $G_1 = G_2$  the derivative vanishes. So, this CSF is not regular. Define  $U_1 =$   
 230  $G_1 + \sqrt{(-\ln x)/(1 - \ln x)} - a$ , where  $(x, a) = (2\theta, 0)$  if  $0 < \theta \leq 1/2$  and  
 231  $(x, a) = (2\theta - 1, 1)$  if  $1/2 < \theta \leq 1$ .<sup>6</sup> Let  $U_2 = G_2$ . Notice that SC holds. Straight-  
 232 forward manipulations show that this pair of utility functions rationalizes the smooth  
 233 difference-form contest in (11).

## 234 2.2 More than two contestants

235 In the case of three contestants, the previous argument does not yield microfounda-  
 236 tions for the additive CSF (5). There are two reasons for that, which are explained  
 237 in Propositions 2 and 3 below. The first result shows that it might be impossible to

<sup>6</sup> One might also define  $U_1 = G_1 + 1$ , when  $\theta = 0$ .

238 partition  $\Theta$  in  $n$  nonempty intervals, which is what is implied by the SC assumption.  
 239 The second result shows that even if such a partition is assumed, the win probability  
 240 of a given contestant might not be responsive to changes in the efforts of all other  
 241 contestants, as in (5). First, we need the following assumption:

242 **Assumption 2**  $U_i(\theta, G_i)$  are continuous and  $U_i(\theta, G_i) \rightarrow \infty$  when  $G_i \rightarrow \infty$ ,  
 243  $i = 1, 2, \dots, n$ .

244 This assumption (A.2 in the sequel) is fulfilled in the payoff functions used in  
 245 Examples 1 and 2 before. In the case of Examples 3 and 4, this assumption is fulfilled  
 246 if  $f_i(G_i) \rightarrow \infty$  when  $G_i \rightarrow \infty$ , which is the case in (1). Thus, it looks like a pretty  
 247 harmless assumption. However, its consequences are not.

248 **Proposition 2** Under Assumption A.2, and when  $n = 3$ , the additive CSF (5) cannot  
 249 be obtained from payoff maximization when SC holds for players 1 and 2.

250 *Proof* Let  $U'_3(G_3) = \max U_3(\theta, G_3), \theta \in \Theta$ . The maximum exists and varies contin-  
 251 uously with  $G_3$  (by Berge's maximum theorem). By taking  $G_1$  and  $G_2$  large enough,  
 252 say  $G'_1$  and  $G'_2$ , the property (SC) and A.2 imply that there is a  $\bar{\theta}$ , such that

$$253 \quad U_1(\theta, G'_1) > U'_3(G_3), \quad \forall \theta \in [0, \bar{\theta})$$

$$254 \quad U_2(\theta, G'_2) > U'_3(G_3), \quad \forall \theta \in (\bar{\theta}, 1].$$

255 Thus, player 3 never obtains the prize. Moreover, because  $U'_3(\cdot)$  is continuous in  $G_3$ ,  
 256 small variations in  $G_3$  do not affect neither  $p_1$  nor  $p_2$ , thus the result.  $\square$

257 Similar results can be obtained for  $n > 3$  by extending suitably the SC condition.  
 258 However, as the next result shows, even weak generalizations of the SC condition  
 259 cause lack of rationalizability of the additive CSF (5) even if Assumption A.2 is not  
 260 postulated. First let us consider the following generalization of SC.

261 **Definition 3** A collection of payoff functions  $U_i(\theta, G_i) i = 1, 2, \dots, n$  satisfies the  
 262 Generalized Single Crossing (GSC) condition when for all  $\mathbf{G}$ , there is a permutation  
 263 in the set of agents  $i, j, \dots, k$  and a partition of  $\Theta$ ,  $(\Theta_i, \Theta_{ij}, \Theta_j, \dots, \Theta_r, \Theta_{rk}, \Theta_k)$   
 264 such that  $\Theta_s = \{\theta \mid U_s(\theta, G_s) > U_r(\theta, G_r), \forall r \neq s\}, s = i, j, \dots, k, \Theta_{sh} =$   
 265  $\{\theta \mid U_s(\theta, G_s) = U_h(\theta, G_h)\}$ , with all  $\Theta_{sh}$  singletons for  $s, h = i, j, \dots, k$ .

266 Notice that, when  $n = 2$ , GSC is implied by SC.

267 **Proposition 3** When the utility functions satisfy the GSC and are continuous, the  
 268 additive CSF (5) cannot be obtained from payoff maximization.

269 *Proof* We will prove the result for  $n = 3$ . The extension to  $n > 3$  is straightforward.  
 270 Without loss of generality, let the permutation of  $N$  be 1, 2, 3. Then,

$$271 \quad U_1(\theta, G_1) > U_j(\theta, G_j), \quad j = 2, 3, \quad \forall \theta \in \Theta_1$$

$$272 \quad U_2(\theta, G_2) > U_j(\theta, G_j), \quad j = 1, 3, \quad \forall \theta \in \Theta_2$$

$$273 \quad U_3(\theta, G_3) > U_j(\theta, G_j), \quad j = 1, 2, \quad \forall \theta \in \Theta_3.$$

274 Thus,  $p_1 = \text{length } \Theta_1$ ,  $p_2 = \text{length } \Theta_2$  and  $p_3 = \text{length } \Theta_3$ . It is clear that  $p_1$  (resp.  
275  $p_3$ ) does not depend on  $G_3$  (resp.  $G_1$ ) for small variations of this variable. Thus, the  
276 required functional form cannot be obtained in this case.  $\square$

277 Notice that the results in Propositions 2 and 3 do not depend on  $F(\theta)$  being uniform.  
278 The reason is that given an interval  $[a, b]$  different distributions assign different prob-  
279 ability mass  $F(b) - F(a)$ . However, in these results, it is crucial that the delimiters  
280  $a$  and  $b$  do not depend on the effort of one contestant. When there are two agents,  
281 delimiters depend on both contestants, because each agent competes with the other,  
282 but when there are three or more agents, some agents may compete with a subset of  
283 other agents and not with all of them.

284 Albeit this difficulty in deriving the additive CSF (5) for more than three contes-  
285 tants, contestants' uncertainty about the type of the contest administrator seems to  
286 be a reasonable approach to CSFs. Therefore, it is an important research program  
287 to find CSFs that are rationalizable according to Definition 2 and to work out the  
288 consequences of these new functional forms on equilibrium, comparative statics, etc.  
289 We show now that although this route appears to be promising, it is not free from  
290 difficulties. We will work out two examples and we will show that in both cases, the  
291 following holds<sup>7</sup>

- 292 • CSFs are neither differentiable nor concave.
- 293 • Despite the symmetric nature of basic data, no symmetric Nash equilibrium exists.

294 *Example 6* Let  $U_1(\theta, G_1) = (1 - \theta)G_1$ ,  $U_2(\theta, G_2) = G_2 2/3$  and  $U_3(\theta, G_3) = \theta G_3$ .  
295 Notice that if  $G_1 = G_2 = G_3$ ,  $p_1 = p_2 = p_3 = 1/3$ . We will compute the best reply  
296 of contestant 1.

297 If  $G_2 2/3 < G_3$ , we have two cases: first, if  $G_1 < G_2 2/3$ , then  $p_1 = 0$ ; second, if  
298  $G_1 \geq G_2 2/3$ , then

$$299 \quad p_1 = \begin{cases} (G_1 - G_2 2/3) / G_1 & \text{if } G_1 < (G_3 G_2 2/3) / (G_3 - G_2 2/3) \\ G_1 / (G_1 + G_3) & \text{otherwise.} \end{cases}$$

300 If  $G_2 2/3 \geq G_3$ , we again have two cases

$$301 \quad p_1 = \begin{cases} 0 & \text{if } G_1 < G_2 2/3 \\ (G_1 - G_2 2/3) / G_1 & \text{otherwise.} \end{cases}$$

302 In a symmetric equilibrium  $\hat{G}$  we have  $G_1 \geq G_2 2/3$  and  $G_1 < (G_3 G_2 2/3) /$   
303  $(G_1 - G_2 2/3)$ . Thus, contestant 1 maximizes  $V(G_1 - G_2 2/3) / G_1 - G_1$ , where  $V$   
304 is the value of the prize. If the equilibrium is symmetric, it must be at positive level  
305 of effort. Thus, the maximum is interior and the first-order condition yields the best  
306 reply, namely  $G_1 = (V G_2 2/3)^{1/2}$ .

307 For  $\hat{G}_1 = \hat{G}_2$ , this yields  $\hat{G}_1 = V 2/3$ . We now have to make sure that this payoff is  
308 larger than the payoff associated to  $G_1 = 0$  (yielding a  $p_1$  and a payoff equal to 0).  
309 This is equivalent to  $\hat{G}_2 \leq V 27/100$ , which contradicts  $\hat{G}_1 = \hat{G}_2 = V 2/3$ .

<sup>7</sup> This may also happen for  $n = 2$  (see Che and Gale 2000).

310 Example 6 can be criticized, because the existence of endpoints (0 and 1) makes  
 311 contestants nonsymmetric. For instance, if  $G_1 = G_2 = G_3$ , a variation of  $G_2$  affects  
 312  $p_1$  and  $p_3$ , but a variation of  $G_1$  only affects  $p_2$ . Thus, we now adapt the model  
 313 of Salop (1979) of a circular city to our framework. Here, symmetry of the effects  
 314 of efforts is restored, since each contestant affects the win probability of all other  
 315 contestants.

316 *Example 7* Suppose that three contestants are symmetrically distributed at locations  
 317  $(l_1, l_2, l_3) = (0, 1/3, 2/3)$  on the unit circle, which is now our set of states of the  
 318 world. Assume that  $U_i(\theta, G_i) = u - k |l_i - \theta| + G_i^\alpha$ , where  $u, k$  and  $\alpha$  are posi-  
 319 tive scalars and  $\alpha \leq 1$ . Notice that when effort levels are similar, the relevant  
 320 competition is pairwise: 1 competes only with 2 (resp. 3) for  $\theta \in [0, 1/3]$  (resp.  
 321  $\theta \in [2/3, 1]$ ), while only 2 and 3 compete for  $\theta \in [1/3, 2/3]$ . Thus, the state of the  
 322 world for which, given efforts, the decider is indifferent between candidates 1 and 2  
 323 is

$$324 \theta_{12} = \frac{1}{6} + \frac{1}{2k} (G_1^\alpha - G_2^\alpha).$$

325 A similar reasoning in the case of 1 and 3 yields

$$326 \theta_{13} = \frac{5}{6} + \frac{1}{2k} (G_3^\alpha - G_1^\alpha).$$

327 This implies that  $p_1 = \theta_{12} + 1 - \theta_{13}$ . To determine the CSF in general, suppose  
 328 without loss of generality that  $G_1 \geq G_2 \geq G_3$ . If  $G_1^\alpha - G_3^\alpha \geq k/3$ , then we obtain a  
 329 generalized version of Che and Gale's two-player contest [given in (7)]

$$330 p_1 = \min \left\{ \frac{1}{2} + \frac{1}{k} (G_1^\alpha - G_2^\alpha), 1 \right\}, p_2 = 1 - p_1 \quad \text{and} \quad p_3 = 0;$$

331 and otherwise

$$332 p_i = \frac{1}{3} + \frac{1}{2k} \left( 2G_i^\alpha - G_j^\alpha - G_k^\alpha \right), \quad \text{for } i = 1, 2, 3 \quad \text{and} \quad i \neq j, k.$$

333 Assume  $\alpha < 1$ . A symmetric equilibrium  $\hat{G}$  requires that  $\hat{G}_1$  maximizes 1's payoffs,  
 334 given  $\hat{G}_2$  and  $\hat{G}_3$  and that  $\hat{G}_1 = \hat{G}_2 = \hat{G}_3$ . Thus,  $\hat{G}_1$  maximizes  $p_1 V - G_1$ , where  
 335  $V$  is the value of the prize. If the maximum is interior,  $\hat{G}_1 = (\alpha V/k)^{1/(1-\alpha)}$ . Thus,  
 336 if payoffs of 1 for this value of efforts are negative, 0 effort is the best reply and no  
 337 symmetric equilibrium exists.

338 Note that it is straightforward to extend the last example to more than three con-  
 339 testants. The so derived CSF can be seen as an extension of Che and Gale's linear  
 340 difference-form [given in (7)] to more than two contestants [see (9)].



## 341 2.3 An alternative notion of rationalizability

342 The simple setting considered so far might be adapted in several ways to yield the  
 343 additive CSF (5) when there are more than three contestants: (i) The type of the  
 344 contest administrator might be multidimensional; (ii) the distribution function might  
 345 be nonuniform; (iii) the rationalizability notion might be different. Given that (i)  
 346 and (ii) have already been explored (e.g. in Hillman and Riley 1989), we pursue now  
 347 (iii).

348 Consider a situation where a contest administrator cares not only about the effort  
 349 of the winner of the contest but also about the effort of others. One might think of  
 350 the promotion of workers in a firm based on their performance or of firms competing  
 351 for a research prize based on R&D investment, which generates new knowledge. In  
 352 such a situation, the type of the decider represents how much he values the effort of a  
 353 particular contestant relative to the others. We present an example yielding a special  
 354 case of the additive CSF (5) for three contestants. This example can easily be extended  
 355 to more agents and to more general effectivity functions.

356 *Example 8* Let  $U_1 = (1 - \theta)G_1 - \theta(G_2 + G_3)$ ,  $U_3 = \theta G_3 - (1 - \theta)(G_1 + G_2)$  and  
 357 normalize  $U_2 = 0$ . We have that

$$358 \quad U_1 \geq U_2 \Leftrightarrow \theta \leq \theta_{12} \equiv \frac{G_1}{G_1 + G_2 + G_3},$$

$$359 \quad U_1 \geq U_3 \Leftrightarrow \theta \leq \theta_{13} \equiv \frac{2G_1 + G_2}{2(G_1 + G_2 + G_3)},$$

$$360 \quad U_3 \geq U_2 \Leftrightarrow \theta \geq \theta_{23} \equiv \frac{G_1 + G_2}{G_1 + G_2 + G_3}.$$

361 This yields

$$362 \quad p_1 = \theta_{12} = \frac{G_1}{G_1 + G_2 + G_3},$$

$$363 \quad p_2 = \theta_{23} - \theta_{12} = \frac{G_2}{G_1 + G_2 + G_3},$$

$$364 \quad p_3 = 1 - \theta_{23} = \frac{G_3}{G_1 + G_2 + G_3}.$$

365 **3 CSFs as sharing rules**

366 Inspired by the second interpretation of the outcome of a contest as shares of the  
 367 prize, we establish now a connection to bargaining and claims problems. This can be  
 368 interpreted as contestants bargaining over all possible assignments of win probabili-  
 369 ties or over shares. If no agreement is reached, all win probabilities are zero. In our  
 370 approach, a variation in effort only affects the share of the prize. A more complete  
 371 theory might consider that the size of the prize is also affected. This allows taking  
 372 into account the opportunity cost of effort (see Anbarci et al. 2002; Garfinkel and  
 373 Skaperdas 2007).

374 3.1 ‘Classical’ bargaining

375 A contest problem is a vector  $f(\mathbf{G}) = (f_1(G_1), \dots, f_n(G_n))$  with at least two entries,  
 376 each of which is strictly positive.<sup>8</sup> Since we consider a fixed vector of efforts  $\mathbf{G}$ , we  
 377 will simply use the notation  $f_i$  instead of  $f_i(G_i)$  and  $f$  instead of  $f(\mathbf{G})$ . An allocation  
 378 in a contest problem is a  $n$ -tuple  $\mathbf{p} = (p_1, \dots, p_n) \in \mathbb{R}^n$  with  $0 \leq p_i \leq 1$  and  
 379  $\sum_{i=1}^n p_i = 1$ . A CSF is a function that assigns a unique allocation to each contest  
 380 problem.

381 We define now a bargaining problem associated with each contest problem. A bar-  
 382 gaining problem is a pair  $(S, \mathbf{d})$  where  $S \subset \mathbb{R}^n$  is a compact convex set,  $\mathbf{d} \in S$ , and  
 383 there exists  $\mathbf{s} \in S$  such that  $s_i > d_i, i = 1, \dots, n$ . The set  $S$ , the feasible set, consists  
 384 of all utility vectors attainable by the  $n$  contestants through unanimous agreement.  
 385 The disagreement point  $\mathbf{d}$  is the utility vector obtained if there is no agreement. In our  
 386 context, it seems natural to define

$$387 \quad S = \left\{ \mathbf{p} \in \mathbb{R}^n \mid 0 \leq p_i \leq 1 \text{ and } \sum_{i=1}^n p_i \leq 1 \right\} \text{ and } \mathbf{d} = \mathbf{0}.$$

388 A bargaining solution is a function  $\psi$  assigning to each bargaining problem  $(S, \mathbf{d})$   
 389 a unique element in  $S$ . We are interested in the weighted Nash solution with weights  
 390  $\alpha$ .

391 **Definition 4** Let  $\alpha_i > 0$  for all  $i = 1, \dots, n$ . The  $\alpha$ -asymmetric Nash solution is  
 392 defined as

$$393 \quad \psi^\alpha = \arg \max_{\mathbf{p} \in S} \prod_{i=1}^n (p_i - d_i)^{\alpha_i}.$$

394 In this framework, it is natural that the effort of a contestant determines his  
 395 bargaining position. Suppose that efforts affect the exponents of the weighted Nash  
 396 bargaining solution as defined before. For simplicity, let  $\alpha = f$ . The next result is  
 397 parallel to one obtained by Dagan and Volij (1993) in a different framework.<sup>9</sup>

398 **Proposition 4** The  $\alpha$ -asymmetric Nash solution for  $\alpha = f$  induces the additive CSF  
 399 (5).

400 *Proof* Let  $f$  be a contest problem; consider the associated bargaining problem and  
 401 let  $\psi^\alpha = \mathbf{p}^*$ . The first-order conditions of the maximization problem defining the  
 402 asymmetric Nash solution with  $\mathbf{d} = \mathbf{0}$  imply that

$$403 \quad p_j^* = \frac{\alpha_j}{\alpha_i} p_i^*, \text{ for all } i, j \in N.$$

<sup>8</sup> If  $f_i(G_i) = 0$  for some contestant  $i$ , assign zero win probability to this agent and consider the reduced vector in which the entry corresponding to agent  $i$  is missing.

<sup>9</sup> In the literature, the weighted Nash solution has also been interpreted as a decider maximizing a payoff function. This is another example of the connections between the approaches taken in Sect. 2 and here.

404 Given the Pareto optimality of the asymmetric Nash solution, we have that  $\sum_{j=1}^n$   
 405  $p_j = 1$ . This implies  $p_i^* = \alpha_i / \sum_{j=1}^n \alpha_j$ .  $\square$

406 Since the preceding result sheds light on the additive CSF (5) from a very different  
 407 angle than the approach of the previous section, it is of interest on its own right. How-  
 408 ever, it also opens the door to understand CSFs as the outcome of strategic bargaining  
 409 models based on Rubinstein's alternating offers game. Since it is well known that  
 410 under certain conditions the asymmetric Nash solution can be supported by such a  
 411 game, it follows that alternative conditions thought to reflect reasonable properties of  
 412 underlying institutional details can yield alternative CSFs.

### 413 3.2 Bargaining with claims

414 It might seem odd that, while the effort vector  $f$  defines a contest problem, this  
 415 information is not used in the description of the associated bargaining problem  $(S, d)$ .  
 416 If we want to incorporate this information in the description of the problem, the rel-  
 417 evant framework is the one of bargaining problems with claims (Chun and Thomson  
 418 1992).<sup>10</sup> A contest-bargaining problem is then a triple  $(S, d, f)$  with the following  
 419 interpretation: contestants bargain over all possible assignments of win probabilities.  
 420 The contestants' effectivity functions translate individual effort into an 'aspiration  
 421 point'  $f$ . Thus,  $f(G)$  measures the social merit that society or the decider awards to  
 422 the vector of efforts  $G$ .

423 If no unanimous agreement is reached, all win probabilities are zero. A contest-  
 424 bargaining solution  $\phi$  assigns to each such triple a unique element in  $S$ . A maximal  
 425 point  $p$  of  $S$  is a point such that  $\sum_{j=1}^n p_j = 1$ . The proportional solution is defined  
 426 as follows.

427 **Definition 5** The proportional solution  $\phi^P$  is defined as the maximal point  $p$  of  $S$  on  
 428 the segment connecting the disagreement point  $d$  and the aspiration point  $f$ .

429 **Proposition 5** *The proportional solution induces the additive CSF (5).*

430 *Proof* Let  $f$  be a contest problem; consider the associated bargaining problem with  
 431 claims and let  $\phi^P = p^*$ . The line that passes through the two points  $d$  and  $f$  is  
 432 the set of vectors  $x$  of the form  $x = (1 - t)d + t f$ , with  $t \in \mathbb{R}$ . Since  $d = 0$ ,  
 433  $x = t f$ . Given that  $p^*$  is a maximal point, we have that  $t = 1 / \sum_{j=1}^n f_j$ . This  
 434 implies  $p_i^* = f_i / \sum_{j=1}^n f_j$ .  $\square$

435 The richer description of bargaining problems with claims has allowed to define an  
 436 alternative solution that also explicitly builds on the aspiration point  $f$ . Bossert (1993)  
 437 analyzes the claim-egalitarian solution. For the purpose of the next proposition, it suf-  
 438 fices to consider the case of two contestants. The following definition is adapted to  
 439 our context, because in contest problems there is no upper bound on individual effort  
 440 levels, that is,  $f$ .

<sup>10</sup> Notice that a contest problem is not equivalent to a bargaining problem with claims. One important difference is that in contest problems there is no upper bound on individual effort levels, that is,  $f$ .

441 **Definition 6** Let  $n = 2$  and  $f_h \geq f_l, h, l = 1, 2$ . The claim-egalitarian solution  $\phi^E$   
 442 is defined as the maximal point  $\mathbf{p}$  of  $S$  such that  $f_h - p_h = f_l - p_l$  if  $f_h - f_l \leq 1$ .  
 443 Otherwise  $p_h = 1$  and  $p_l = 0$ .

444 The claim-egalitarian solution selects a point on the Pareto frontier of  $S$  such that  
 445 the loss of each contestant compared with his aspiration level is the same for all agents  
 446 (if such a point exists). This is an egalitarian solution in the sense that the absolute  
 447 amount each agent has to give up is equalized across contestants. The next proposition  
 448 says that this idea is the same as saying that only differences in effort matter.

449 **Proposition 6** For  $n = 2$ , the claim-egalitarian solution induces a generalization of  
 450 Che and Gale's difference-form CSF, that is,

$$451 \quad \phi_i^E = p_i^{CG'}(\mathbf{G}) = \max \left\{ \min \left\{ \frac{1}{2} + \frac{1}{2} (f_i - f_j), 1 \right\}, 0 \right\} \quad \text{for } i = 1, 2.$$

452 *Proof* The fact that if  $|f_i - f_j| \geq 1$  then  $\phi_i^E = p_i^{CG'}(\mathbf{G})$  is obvious. Suppose  
 453  $|f_i - f_j| \leq 1$ . Since  $p_j = 1 - p_i$ , we have  $f_i - p_i = f_j - (1 - p_i)$ . Rearranging  
 454 yields the desired expression.  $\square$

455 Notice that when  $f_i(G_i) = 2\sigma G_i$ , for  $i = 1, 2$ , where  $\sigma$  is a positive scalar, we  
 456 obtain (7), the class of linear difference-form functions analyzed by Che and Gale  
 457 (2000). Notice that it is straightforward to extend the last result to more than two con-  
 458 testants [see (9)].<sup>11</sup> Interestingly, this recommendation differs in the minimal effort  
 459 necessary to obtain a nonzero share and in the marginal effect of effort from the one  
 460 based on Example 7.

461 Definition 6 equalizes losses based on absolute claims. This creates the 'kink' and  
 462 the nonresponsiveness of Che and Gale's CSF to effort when the difference in aspi-  
 463 ration levels is high enough. Considering relative claims, this can be avoided. Notice  
 464 that  $f_i/f_h$  (for  $i = 1, \dots, n$ ) indicates the percentage contestant  $i$ 's aspiration level  
 465  $f_i$  constitutes of the highest level  $f_h$ .

466 **Definition 7** Let  $n = 2$  and w.l.o.g. denote  $f_h = \max\{f_1, f_2\}$ . The relative  
 467 claim-egalitarian solution  $\phi^{RE}$  is defined as the maximal point  $\mathbf{p}$  of  $S$  such that  
 468  $f_1/f_h - p_1 = f_2/f_h - p_2$ .

469 The relative claim-egalitarian solution selects a point on the Pareto frontier of  $S$   
 470 such that the loss of each contestant compared with this 'relative claim point' is the  
 471 same for all agents. The next proposition relates this idea to the serial CSF.<sup>12</sup>

<sup>11</sup> For  $n = 3$  and  $f_1 \geq f_2 \geq f_3$ , it is natural to require the following. If  $f_1 - f_2 \geq 1$ , then  $p_1 = 1$  and  $p_2 = p_3 = 0$ . If  $f_1 - f_3 \geq 1 > f_1 - f_2$ , then  $\phi^E$  is the maximal point  $\mathbf{p}$  of  $S$  such that  $p_3 = 0$  and  $f_1 - p_1 = f_2 - p_2$ . Lastly, when  $f_1 - f_3 < 1$ , then  $\phi^E$  is the maximal point  $\mathbf{p}$  of  $S$  such that  $f_1 - p_1 = f_2 - p_2 = f_3 - p_3$ .

<sup>12</sup> This reasoning can easily be extended to more contestants. However, the requirement that  $f_i/f_h - p_i = f_{i+1}/f_h - p_{i+1}$  for all  $i = 1, \dots, n - 1$  does not always yield welldefined win probabilities. A way out is the following. Consider an ordered vector  $f_1 \geq f_2 \geq \dots \geq f_n$  and rescale the 'relative claim point' to make the pairwise comparisons  $f_i/(i \cdot f_h) - p_i = f_{i+1}/(i \cdot f_h) - p_{i+1}$  for all  $i = 1, \dots, n - 1$ . This coincides with Definition 7 when there are two agents and yields a generalization of the serial CSF for any number of contestants.

472 **Proposition 7** For  $n = 2$  and  $f_1 \geq f_2$ , the relative claim-egalitarian solution induces  
 473 a generalization of the serial CSF, that is,

$$474 \quad \phi_i^{\text{RE}} = p_i^{S'}(\mathbf{G}) = \sum_{j=i}^2 \frac{f_j - f_{j+1}}{j \cdot f_h} \quad \text{for } i = 1, 2 \quad \text{and } f_3 = 0.$$

475 *Proof* W.l.o.g. assume  $f_1 \geq f_2$ . We have that  $1 - p_1 = f_2/f_1 - p_2 = f_2/f_1 - 1 + p_1$ .  
 476 This can be rewritten as  $p_1 = 1 - f_2/(2f_1) = (f_1 - f_2)/f_1 + f_2/(2f_1)$ . Since  $\phi^{\text{RE}}$   
 477 must be a maximal point, we obtain  $p_2 = f_2/(2f_1)$ .  $\square$

#### 478 4 Concluding remarks

479 In line with two prominent interpretations of the outcome of contests, this paper has  
 480 investigated foundations for prominent CSFs based on two different approaches. The  
 481 first analyzes the implications of contestants' incomplete information concerning the  
 482 'type' of the contest administrator. The second understands CSFs as sharing rules  
 483 and makes a connection to bargaining and claims problems. Both approaches provide  
 484 foundations for popular CSFs and guidelines for the definition of new ones. The results  
 485 of this paper suggest two lines for future research on CSFs.

486 On the normative side, the implications of linking the problem of assigning win  
 487 probabilities in contests to bargaining, claims and taxation problems are twofold.

488 It might yield an improved understanding of existing CSFs. For instance, pro-  
 489 portionality principles have been defended at least since the philosophers of ancient  
 490 Greece. Therefore, it seems possible to obtain different characterizations of the addi-  
 491 tive CSF (5) using ideas of characterizations of proportionality stressed in these related  
 492 problems.<sup>13</sup>

493 It suggests guidelines for the definition of new CSFs, since different normative prin-  
 494 ciples might lead to the formulation of different classes of CSFs. A case in point here  
 495 is the claim-egalitarian solution that gives a recommendation as to how to extend the  
 496 difference-form functions analyzed by Che and Gale (2000) to more than two contes-  
 497 tants. On the positive side, the implications for future research parallel the normative  
 498 ones.

499 Solution concepts in bargaining, claims and taxation problems that can be related to  
 500 popular CSFs might yield rationales for the latter. An example is to link contests with  
 501 the Bilateral Principle that has proved a fruitful way to incorporate Luce's Choice  
 502 Axiom into game theory. Dagan et al. (1997) have provided a game form captur-  
 503 ing the non-cooperative dimension of the consistency property of bankruptcy rules.<sup>14</sup>  
 504 An adaptation of their result in our framework shows that the additive CSF (5) can be

<sup>13</sup> Note that the class of problems in which win probabilities are assigned has a particularly simple structure. This implies that a characterization of a solution for a larger class of problems does not need to characterize a solution for contests.

<sup>14</sup> Notice that a contest problem is not equivalent to a bankruptcy problem in which the estate is equal to one, since in contest problems there is no lower bound on the sum of individual effort levels, that is,  $\sum_{j=1}^n f_j$ .

505 supported by a pure strategy subgame perfect equilibrium of a certain non-cooperative  
506 game.

507 By incorporating realistic details of contest situations, novel CSFs can be derived.  
508 Examples are the recommendation of the circular model in Example 7 as to how to  
509 extend Che and Gale's difference-form function to more than two contestants or the  
510 effects of modifying Rubinstein's alternating offers bargaining game.

511 Lastly, we remark that there is no straightforward generalization of the single-crossing  
512 property that would generalize the results of Sect. 2.1 to more than two contestants.  
513 In any case, there might be a way to allow for conditional preferences over subsets of  
514 players (e.g. in the spirit of Luce's Choice Axiom) that would allow for a representation  
515 theorem. We leave this to future research.

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